

**EUR 3559 e**

**EUROPEAN ATOMIC ENERGY COMMUNITY - EURATOM**

**AN APACHE "SUBROUTINE" FOR GENERATING PADE'S CIRCUITS**

**by**

**C. BONA**

**1967**



**Joint Nuclear Research Center  
Ispra Establishment - Italy**

**Scientific Information Processing Center - CETIS**



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## **SUMMARY**

With APACHE it is possible to write programs corresponding to analog circuits of common use which can be considered as the "Subroutines" or "Functions" in digital programs.

In fact, when a problem requires a circuit for which a subroutine has been written one can avoid to write the corresponding equations just by adding the deck constituting the subroutine to the APACHE deck representing the problem.

In the following, two subroutines for generating any number of Pade's circuits for fixed and variable time delays are described.

## **KEYWORDS**

PROGRAMMING,  
COMPUTERS,  
ANALOG SYSTEMS,  
DIGITAL SYSTEMS,  
CIRCUITS.

## **Codes**

## An APACHE "Subroutine" for Generating Pade's circuits (+)

### INTRODUCTION

When, in the framework of a larger job, a circuit to simulate a fixed or variable delay is needed, a lot of tedious work has to be done to implement one of the many circuits which approximate the delay operator (Pade approximation, Bode approximation, etc.)

In fact there are no conceptual difficulties, as these circuits are well defined and very efficient formulas are provided to calculate the order of approximation needed based on the maximum frequency to be delayed and on the length of the delay.

The only difficulty which may arise is the choice of the scale factors to be given to each intermediate variable, the other things that one has to do being routine works like computing the pot setting values and drawing the circuit etc.

With APACHE (foot note 1) all these routine works are avoided, being automatically done by the compiler, while a new problem arises: the given transfer functions must be transformed in a system of differential equations of first order; besides an estimate must be given of the maximum values of all variables.

On the other side with APACHE the program corresponding to a given circuit can be written once for ever, punched on cards and kept for every time it is needed.

But there is still something more, with APACHE one need not to write for each order of approximation the

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(+) Manuscript received on May 22, 1967.

appropriate program, it is possible to write the program in such a way to generate as many different delay circuits as needed , and each of them of the desired order.

This has been done for the Pade circuit both for fixed and variable time delays.

The aim of this work is to present these two APACHE "Subroutines" (foot note 1) and to give an idea of how they have been written.

#### How to use the programs

The procedure for using these programs, which consist of two decks of approximately 50 cards each, is very simple:

- 1) Add the deck (or the decks) to the deck of cards representing your problem
- 2) Add the following statements: (only those requested by the type of delay you need in your problem)
  - a) For fixed time delays

#### EQUATIONS

FF(i) = name of the  $i^{\text{th}}$  variable to be delayed  
 .  
 .  
 .  
 Name of the  $i^{\text{th}}$  variable delayed = GF(i)

#### PARAMETERS

NRITF = number of delay circuits requested  
 .  
 .  
 .  
 ORITF(i) = order of the  $i^{\text{th}}$  delay circuit (up to 20)  
 .  
 .  
 .  
 TRFIX(i) = value (in sec.) of the  $i^{\text{th}}$  delay  
 .  
 .  
 .  
 FFINZ(i) = value of the  $i^{\text{th}}$  variable to be delayed at  $t=0$   
 .  
 .  
 .  
 FFMAX(i) = maximum value of the  $i^{\text{th}}$  variable to be delayed



## b) For variable time delays

EQUATIONS

FV(i) = name of the  $i^{\text{th}}$  variable to be delayed  
 ...  
 Name of the  $i^{\text{th}}$  variable delayed = GV(i)  
 ...

PARAMETERS

NRITV = number of delay circuits requested  
 ...  
 ORITV(i) = order of the  $i^{\text{th}}$  delay circuit (up to 20)  
 ...  
 TVINZ(i) = length of the  $i^{\text{th}}$  delay at  $t=0$   
 ...  
 TVMIN(i) = maximum length of the  $i^{\text{th}}$  delay  
 ...  
 TVMAX(i) = maximum length of the  $i^{\text{th}}$  delay  
 ...  
 TVINZ(i) = value of the  $i^{\text{th}}$  variable to be delayed at  $t=0$   
 ...  
 FVMAX(i) = maximum value of the  $i^{\text{th}}$  variable to be delayed

Pade's approximation of  $e^{-p\tau}$ 

As is well known, the transfer function of a pure time delay can be approximated as follows ( $n^{\text{th}}$  order Pade's approximation):

$$1) \quad G(s)/F(s) = e^{-s\tau} \approx P_n/Q_n$$

where  $G(s) = \mathcal{L}(F(t-\tau))$  and  $F(s) = \mathcal{L}(F(t))$  and

$$P_n = h_0 - h_1 \cdot (s\tau) + \dots + (-1)^i \cdot h_i \cdot (s\tau)^i + \dots + (-1)^n \cdot h_n \cdot (s\tau)^n$$

$$Q_n = h_0 + h_1 \cdot (s\tau) + \dots + h_i \cdot (s\tau)^i + \dots + h_n \cdot (s\tau)^n$$

$$h_0 = 1, h_1 = 1/2, \dots, h_i = \frac{n(n-1)\dots(n-i+1)}{2n(2n-1)\dots(2n-i+1) \cdot i!}, \dots$$

Note that  $P_n$  is quite the same as  $Q_n$  except for the sign of the  $i^{\text{th}}$  terms for odd values of  $i$ .

We can write:  $G(s) \cdot Q_n = F(s) \cdot P_n$  from which, dividing both terms by  $(s\tau)^n$ , we obtain:

$$2) \frac{h_0}{(s\tau)^n} \cdot (F-G) - \frac{h_1}{(s\tau)^{n-1}} \cdot (F+G) + \dots + \frac{(-1)^1 \cdot h_1}{(s\tau)^{n-1}} (F - (-1)^1 \cdot G) + \dots + \\ + \frac{(-1)^n \cdot h_n}{(s\tau)^0} \cdot (F - (-1)^n \cdot G) = 0$$

and substituting  $k_{n-i} \frac{n!}{(2n)!}$  to  $h_i$

$$3) \sum_{i=0}^n \frac{k_i^n}{(s\tau)^i} \left[ (-1)^{n-i} \cdot F - G \right] = 0 \quad k_0 = 1, \dots, k_i^n = \frac{(n+i)!}{i!(n-i)!}$$

which is the classical expression from which the Pade's circuit is deduced.

#### APACHE equations for the Pade's circuit

The easiest way to deduce from eq.3 the Pade's circuit is to rewrite it as follows .

$$3') \quad G = (-1)^n \cdot F + Q_1 \quad ; \quad Q_1 = \sum_{i=1}^n \frac{k_i}{(s\tau)^i} \left[ (-1)^{n-i} \cdot F - G \right]$$

Thus  $G$  can be obtained by adding to  $Q_1 + F$  or  $-F$  depending on the order  $(n)$  of the Pade desired.

$Q_1$  in its turn is made up by integrating several times terms like  $(-F-G)$  and  $(+F-G)$  and by summing them up.

Fig.1 shows, in three steps, for clearness' sake, how  $Q_1$  can be obtained by a simple chain of amplifiers. (no sign inversion is considered here)

Note how the gains have been split between stages; this makes them more uniform and much more easy to be computed as



$$\frac{k_1^n}{k_{i-1}^n} = \frac{(n+1)(n-i+1)}{i}$$

Having  $Q_1$ ,  $G$  is immediately obtained by combining it with  $F$  and the circuit completed by feeding back  $F \pm G$ .

In fig. 2 both schemes for  $n$  odd and even are reported.

The circuit for variable time delays can be deduced in the same way, but it must be remembered that eq.2 is no more rigorous, as well as the splitting of the term  $\frac{1}{\tau^n}$  between the different integration stages.

However, experience shows that the circuits in fig.3 give a rather good approximation of a variable time delay provided that the delay ( $\tau$ ) is slowly varying with respect to  $F$ .

The APACHE equations for the Pade circuit for fixed and variable delays are easily deduced from fig. 2 and 3 where the circuits for an even and an odd Pade order ( $N$ ) are reported. (no sign inversion is considered there)

The names between square brackets are the names used in the APACHE program.

Note that, for reasons which will be explained later, the variables which are output from integrators may have two different names depending on the I.C. value of the function to be delayed.

In fig. 4 and 5 the complete listing of the APACHE programs for fixed and variable Pade circuits are reported.

We leave to the readers to go through the equations in order to see how they correspond to the circuits in fig. 2 and 3.

Note how, by means of the APACHE "DO" statement and of the option "NULL", it has been possible to write one program for generating whatever number of Pade circuits of any desired order (up to 20). (Foot note 2)

### Variables scaling

No problem arises for scaling variables like  $GF(k)$  and  $GV(k)$  (fig. 2 and 3) whose maximum value is obviously equal to the maximum value of the variables  $FF(k)$  and  $FV(k)$  to be delayed, (these values, called  $FFMAX(k)$  and  $FVMAX(k)$ , must be given by the user).

$FQFIX(k)$ ,  $GQFIX(k)$ ,  $QNFIX(k,i)$  which, according to the Pade's order are equal to  $\pm FF(k)$ ,  $-GF(k)$ ,  $\pm FF(k)+GF(k)$ , have been scaled assuming a maximum equal to  $2 * FFMAX(k)$ , which is conservative.

To scale the intermediate variables  $QNFIX(k,i)$  ( $QNVAR(k,i)$  and  $SNVAR(k,i)$  for variable delays), whose maximum values are not easily predictable, the transfer function between each of them and  $FF(k)$  has been calculated and then, looking at their Bode diagrams, it has been found that, for any value of  $k$  and  $i$ , their amplitude never exceeds the value of 2.2, thus even for these variables a maximum value equal to  $2 * FFMAX(k)$  has been taken.

As an example, the amplitude vs. frequency diagrams of  $\mathcal{L}(QNFIX(k,i))/\mathcal{L}(FF(k))$  for a Pade of the fourth order have been reported in fig. 6 and 7.



Set of INITIAL CONDITION values for static check

The IC values of the variables of a Pade circuit depend on the values of the function to be delayed and its derivatives  $\tau$  seconds earlier; we supposed that our Pade circuits will always be used in connection with systems being stationary for  $t \leq 0$ , (foot note 3) thus the IC value for the Pade can be easily deduced imposing that all derivative be zero; we have : (fig. 2)

$$Q_1(t=0) = \begin{cases} N \text{ EVEN} & \begin{cases} 2.F(t=0) & \text{for } i \text{ EVEN} \\ 0 & \text{for } i \text{ ODD} \end{cases} \\ N \text{ ODD} & \begin{cases} 0 & \text{for } i \text{ EVEN} \\ 2.F(t=0) & \text{for } i \text{ ODD} \end{cases} \end{cases}$$

(N = Pade's circuit order)

In order to be able to make a complete static check of the circuit, the I.C. conditions which are zero must be changed into some new fictive IC conditions to be taken from the TEST reference.

The problem, which is well known by analogists, is to choose some fictive IC condition which causes the voltages all over the circuit to be different from zero but not greater than 100V and at the same time to minimize the number of potentiometers requested for it.

In fig. 8 the schemes in fig. 2 are repeated writing beside each element the corresponding IC condition for both cases  $F(t=0) \neq 0$  and  $F(t=0)=0$ .

(This figure corresponds to the analog schemes as deduced from the APACHE output listing)

Note the additional terms from the TEST reference which have been added to the summers which compute GF(k), FQFIX(k) and GQFIX(k) to have a complete set of IC  $\neq$  0 and 100V with only two potentiometers. (Foot note 3)

As it is impossible to use the NULL option in connection with the variable list (i.e. to introduce the comment TEST in a variable definition depending on the value of its IC), to cover all possible cases we have been obliged to write three sets of equations, which hold respectively for:

- 1)  $F(t=0) = 0$
- 2)  $F(t=0) \neq 0$       N EVEN
- 3)  $F(t=0) \neq 0$       N ODD

These three sets of equations are all included in the deck, the appropriate set being selected by deleting the other two by means of the NULL option (see fig. 4 and 5 and coefficient CXF(j), GXF(j) and the DO statement DO, j=ORIFT(k),...)

## CONCLUSIONS

Our experience has shown that to the computer user a lot of thinking and of tedious work is avoided by means of these routines, therefore we are looking forward to see if it is possible to write APACHE subroutines for other circuits of common use. (e.g. for blocks whose transfer function is assigned).

On the other side, the APACHE language, although it has shown to be versatile enough to cover all possible cases for the Padé circuits, appeared rather cumbersome in use for similar problems, because of the lack of "EQUIVALENCE" statements and of



## II

some sort of IF statement (or NULL option) at Parameter and Variable level, which oblige to use fixed names, indexed variables etc. (see above).

Therefore, in the new APACHE version, which is being written now, all these points will be considered.

We hope that these possibilities of organizing the APACHE programs in subroutines will offer still more comfort to the analogue computer user.

Foot note 1

APACHE is a digital program which, starting from a set of equations written in a very intuitive language, performs all the necessary operations to implement them on an analog computer i.e.

a) Reduces the equation in a Standard form corresponding to the characteristic of the analogue elements which will compute them.

b) Performs the scaling of the equations and computes the setting values of potentiometers.

c) Computes the static check value of every element in the circuits.

For further details see : Simulation March 1965 and January 1966 and item 1 in the enclosed bibliography.

An attempt of writing two Apache subroutines like those presented in this paper was done since 1963 by J. Gamp (Bibliography 3 and 4) but only with the last version of the code (July 1965) it has been possible to write them in a really general and usable way.

Obviously, due to the parallel nature of the Apache language, the similarity between APACHE and Fortran subroutines is only formal and consists in the fact that both represent a part of the program that is written once for all.

Thus the user will not find neither CALL or RETURN or similar statements while the physical location of the subroutine deck in the main program is immaterial.

Foot note 2

The APACHE "DO" statement is formally similar to the FORTRAN DO statement but has a different meaning. The APACHE DO is only a means to get greater compactness when many similar equations have to be written (ex. partial differential equations implemented by finite difference method); it does not mean: perform the computation within the DO range for all the indicated values of the indexes, but : consider the equation within the DO range as being written for all the indicated values of the indexes and generate the corresponding circuits.

The NULL option is a means to delete some equations or some part of them: any equation or term which is multiplied by a parameter whose value in the parameter list is 0, NULL is automatically deleted from the problem.

The aim of this option is to give the possibility to include in the range of a DO equations which differ between each other depending on the value of the index.

Foot note 3

The terms  $\pm TREF$ , which in the Apache listing will appear as an "external" variables, correspond to the test reference voltage; the elements which have an input fed by this variable must be connected to  $\pm TEST REF$ .

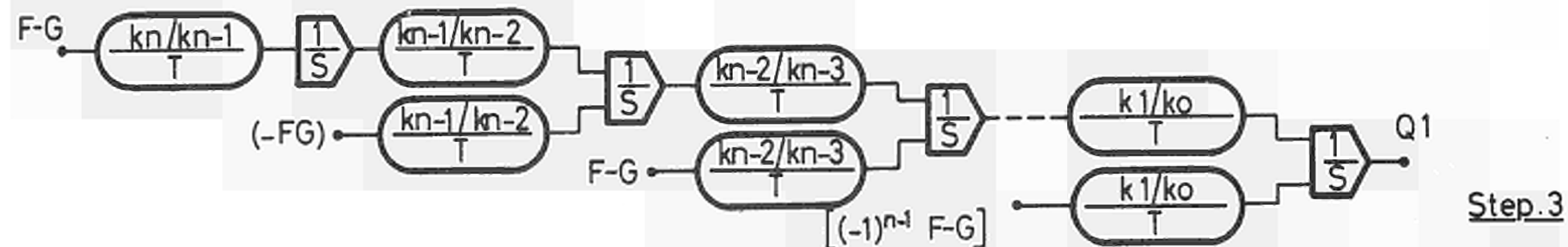
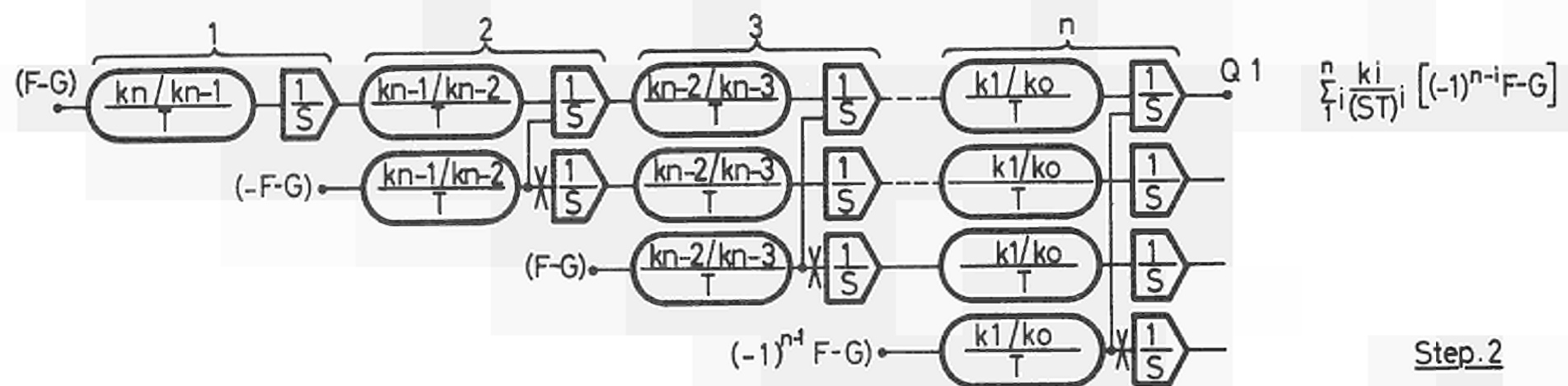
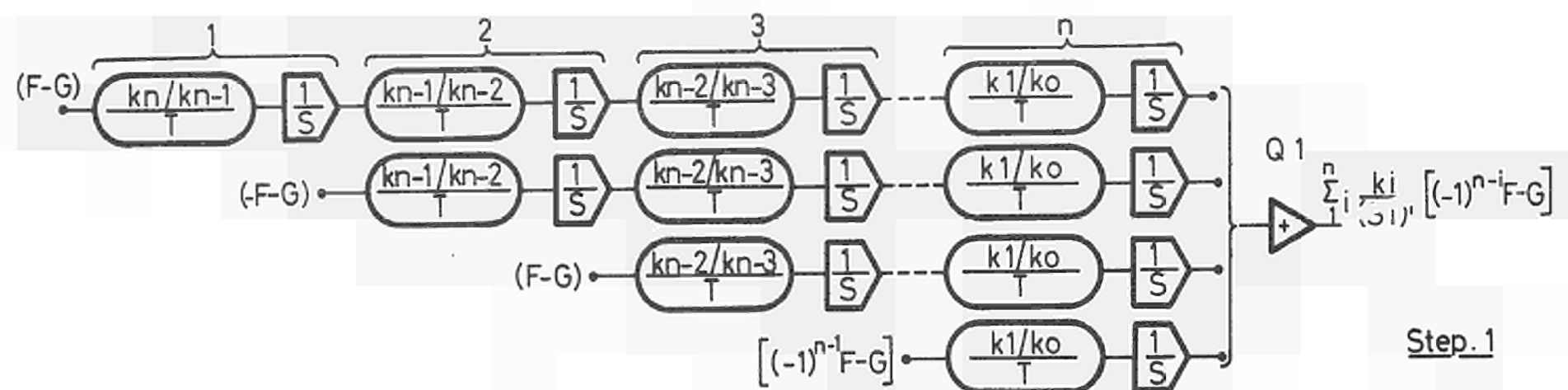


Fig. 1 Implementation of  $Q1 = \sum_{i=1}^n \frac{k_i}{(s\tau)^i} [(-1)^{n-i} \cdot F-G]$ .



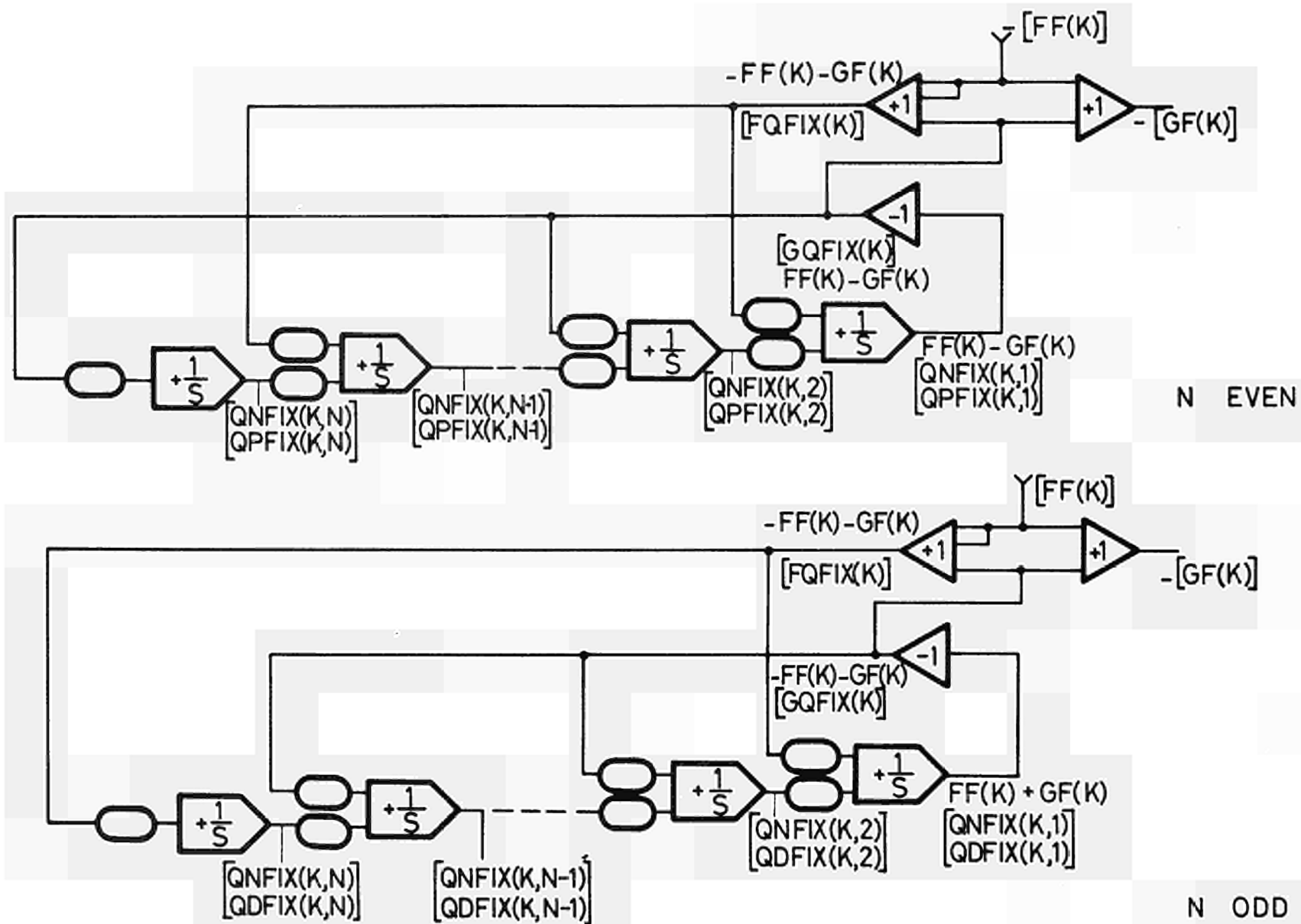


Fig. 2 Pade circuits for fixed time delays.  
Odd and even order approximations (N)

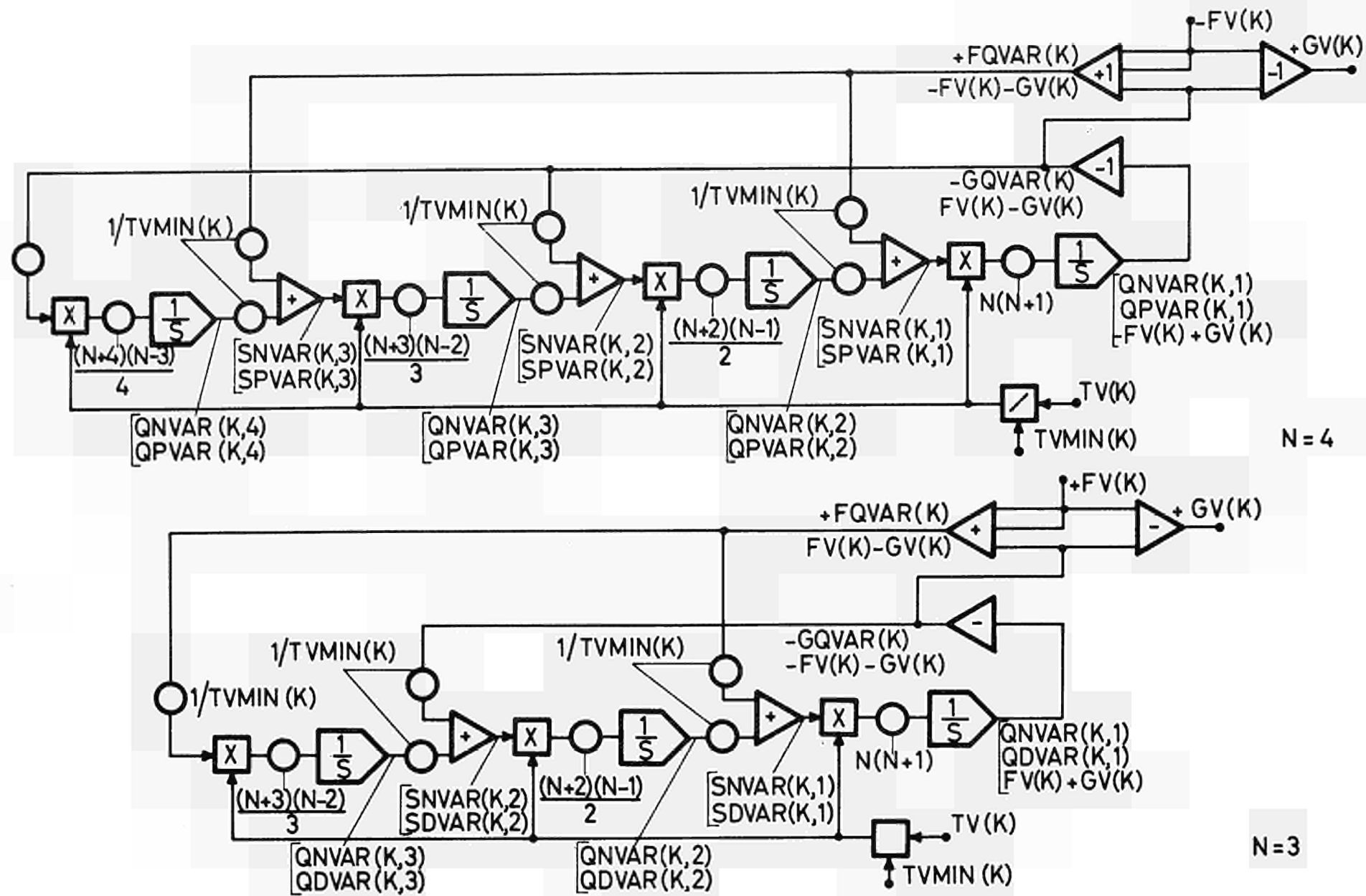


Fig. 3 Pade circuits for variable time delays.  
Odd and even order approximation

PARAMETERS		RF001
	DO,9999,I=1,20,2	RF002
	CDISP(I)=1	RF003
9999	CPARI(I)=0, NULL	RF004
	DO,9998,I=2,20,2	RF005
	CDISP(I)=0, NULL	RF006
9998	CPARI(I)=1	RF007
	CXF(0)=1	RF008
	CXF(1)=0, NULL	RF009
	GXF(0)=0, NULL	RF010
	GXF(1)=1	RF011
	DO,9997,K=1,NRITF	RF012
	MXMF(K)=ORITF(K)+1	RF013
	MXF(K)=ORITF(K)-1	RF014
	LXF(K)=(FFINZ(K)-0)/(FFINZ(K)-0)	RF015
	PCF(K,1)=ORITF(K)*(ORITF(K)+1)	RF016
	DO,9996,I=2,ORITF(K)	RF017
9996	PCF(K,I)=(ORITF(K)+I)*(ORITF(K)-I+1)/I	RF018
9997	QFMAX(K)=2*FFMAX(K)	RF019
VARIABLES		RF020
	TREFF=1,1,TEST	RF020 1
	DO,9995,K=1,NRITF	RF021
	FF(K)=FFINZ(K),FFMAX(K),EXACT	RF022
	GF(K)=,FFMAX(K),EXACT	RF023
	QNFIX(K,1)=0, QFMAX(K),EXACT	RF024
	DO,9994,I=2,ORITF(K)	RF025
9994	QNFIX(K,I)=-QFMAX(K),QFMAX(K),EXACT,TEST	RF026
	QPFIX(K,1)=0, QFMAX(K),EXACT	RF027
	QPFIX(K,2)=-2*FFINZ(K),QFMAX(K),EXACT	RF028
	DO,9993,I=4,MXMF(K),2	RF029
	QPFIX(K,I-1)=-QFMAX(K),QFMAX(K),EXACT,TEST	RF030
9993	QPFIX(K,I)=-2*FFINZ(K),QFMAX(K),EXACT	RF031
	QDFIX(K,1)=-2*FFINZ(K),QFMAX(K),EXACT	RF032
	DO,9992,I=3,MXMF(K),2	RF033
	QDFIX(K,I-1)=-QFMAX(K),QFMAX(K),EXACT,TEST	RF034
9992	QDFIX(K,I)=-2*FFINZ(K),QFMAX(K),EXACT	RF035
	GQFIX(K)=,QFMAX(K),EXACT	RF035 1
9995	FQFIX(K)=,2*FFMAX(K),EXACT	RF036
EQUATIONS		RF037
	DO,9991,K=1,NRITF	RF038
	DO,9991,J=ORITF(K),ORITF(K)	RF039
	DO,9991,S=LXF(K),LXF(K)	RF040
	FQFIX(K)=2*FF(K)+CXF(S)*(GQFIX(K)+1.8*QFMAX(K)*TREFF)+GXF(S)*CPARI	RF041 1
	1(J)*(GQFIX(K)+0.9*QFMAX(K)*TREFF)+GXF(S)*CDISP(J)*(GQFIX(K)+0.9*QF	RF041 2
	2MAX(K)*TREFF)	RF041 3
	GQFIX(K)=CXF(S)*(QNFIX(K,1)-0.9*QFMAX(K)*TREFF)+GXF(S)*CPARI(J)*(Q	RF041 4
	1PFIX(K,1)-0.9*QFMAX(K)*TREFF)+GXF(S)*CDISP(J)*QDFIX(K,1)	RF041 5
	GF(K)=CXF(S)*CPARI(J)*(GQFIX(K)+0.5*QFMAX(K)*TREFF+FF(K))-CXF(S)*C	RF042 1
	1DISP(J)*(GQFIX(K)+0.5*QFMAX(K)*TREFF+FF(K))+GXF(S)*CPARI(J)*(GQFIX	RF042 2
	2(K)+0.9*QFMAX(K)*TREFF+FF(K))-GXF(S)*CDISP(J)*(GQFIX(K)+FF(K))	RF042 3
	DO,9990,I=1,MXF(K)	RF043
9990	CXF(S)*TF(K)*DER1(QNFIX(K,I))=PCF(K,I)*(-QNFIX(K,I+1)+GQFIX(K)	RF044
	1*CPARI(I)-FQFIX(K)*CDISP(I))	RF044
	CXF(S)*CPARI(J)*TF(K)*DER1(QNFIX(K,J))=PCF(K,J)*GQFIX(K)	RF045
	CXF(S)*CDISP(J)*TF(K)*DER1(QNFIX(K,J))=-PCF(K,J)*FQFIX(K)	RF046
	DO,9989,I=1,MXF(K)	RF047
9989	GXF(S)*CPARI(J)*TF(K)*DER1(QPFIX(K,I))=-PCF(K,I)*(QPFIX(K,I+1)+F	RF048
	1QFIX(K)*CDISP(I)-GQFIX(K)*CPARI(I))	RF048
	GXF(S)*CPARI(J)*TF(K)*DER1(QPFIX(K,J))=PCF(K,J)*GQFIX(K)	RF049
	DO,9988,I=1,MXF(K)	RF050
9988	GXF(S)*CDISP(J)*TF(K)*DER1(QDFIX(K,I))=-PCF(K,I)*(QDFIX(K,I+1)-G	RF051
	1QFIX(K)*CDISP(I)+FQFIX(K)*CDISP(I))	RF051
9991	GXF(S)*CDISP(J)*TF(K)*DER1(QDFIX(K,J))=-PCF(K,J)*FQFIX(K)	RF052

Fig. 4 Listing of the APACHE subroutine for fixed time delays.

PARAMETERS		RV000
	DO, 9980, I=1, 20, 2	RV001
	VDISP(I)=1	RV002
9980	VPARI(I)=0, NULL	RV003
	DO, 9979, I=2, 20, 2	RV004
	VDISP(I)=0, NULL	RV005
9979	VPARI(I)=1	RV006
	CXV(0)=1	RV007
	CXV(1)=0, NULL	RV008
	GXV(0)=0, NULL	RV009
	GXV(1)=1	RV010
	DO, 9978, K=1, NRITV	RV011
	MXMV(K)=ORITV(K)-1	RV012
	MXVM(K)=ORITV(K)+1	RV013
	LXV(K)=(FVINZ(K)-0)/(FVINZ(K)-0)	RV014
	PCOV(K, 1)=ORITV(K)*(ORITV(K)+1)	RV015
	DO, 9977, I=2, ORITV(K)	RV016
9977	PCOV(K, I)=(ORITV(K)+1)*(ORITV(K)-I+1)/I	RV017
	QVMAX(K)=2*FVMAX(K)	RV018
9978	SVMAX(K)=2*FVMAX(K)/TVMIN(K)	RV019
VARIABLES		RV020
	TREFV=1, 1, TEST	RV020 1
	DO, 9976, K=1, NRITV	RV021
	FV(K)=FVINZ(K), FVMAX(K), EXACT	RV022
	GV(K)=FVMAX(K), EXACT	RV023
	TV(K)=TVINZ(K), TVMAX(K), EXACT	RV024
	MINRV(K)=TVMAX(K), EXACT	RV025
	UDTXV(K)=1	RV026
	FQVAR(K)=2*FVMAX(K), EXACT	RV027
	GQVAR(K)=QVMAX(K), EXACT	RV027 1
	QNVAR(K, 1)=0, QVMAX(K), EXACT, TEST	RV028
	SNVAR(K, 1)=SVMAX(K), EXACT	RV029
	DO, 9975, I=2, ORITV(K)	RV030
	QNVAR(K, I)=QVMAX(K), QVMAX(K), EXACT, TEST	RV031
9975	SNVAR(K, I)=SVMAX(K), EXACT	RV032
	QPVAR(K, 1)=0, QVMAX(K), EXACT, TEST	RV033
	SPVAR(K, 1)=SVMAX(K), EXACT	RV034
	DO, 9974, I=3, MXVM(K), 2	RV035
	QPVAR(K, I-1)=2*FVINZ(K), QVMAX(K), EXACT	RV036
	SPVAR(K, I-1)=SVMAX(K), EXACT	RV037
	QPVAR(K, I)=QVMAX(K), QVMAX(K), EXACT, TEST	RV038
9974	SPVAR(K, I)=SVMAX(K), EXACT	RV039
	DO, 9976, I=2, MXVM(K), 2	RV040
	QDVAR(K, I-1)=2*FVINZ(K), QVMAX(K), EXACT	RV041
	SDVAR(K, I-1)=SVMAX(K), EXACT	RV042
	QDVAR(K, I)=QVMAX(K), QVMAX(K), EXACT, TEST	RV043
9976	SDVAR(K, I)=SVMAX(K), EXACT	RV044
EQUATIONS		RV045
	DO, 9973, K=1, NRITV	RV046
	DO, 9973, J=ORITV(K), ORITV(K)	RV047
	DO, 9973, S=LXV(K), LXV(K)	RV048
	GV(K)=CXV(S)*(GQVAR(K)-0.5*QVMAX(K)*TREFV)+GXV(S)*VPARI(J)*(GQVAR(K)-0.9*QVMAX(K)*TREFV)+GXV(S)*VDISP(J)*GQVAR(K)+VPARI(J)*FV(K)-VDISP(J)*FV(K)	RV049
	FQVAR(K)=CXV(S)*GQVAR(K)-GXV(S)*VPARI(J)*(GQVAR(K)-0.9*QVMAX(K)*TREFV)-GXV(S)*VDISP(J)*(GQVAR(K)+0.9*QVMAX(K)*TREFV)-VPARI(J)*2*FV(K)+VDISP(J)*2*FV(K)	RV050 1
	GQVAR(K)=CXV(S)*(QNVAR(K, 1)+0.9*QVMAX(K)*TREFV)+GXV(S)*VPARI(J)*(QNVAR(K, 1)+0.9*QVMAX(K)*TREFV)+GXV(S)*VDISP(J)*QDVAR(K, 1)	RV051 1
	MINRV(K)=TVMIN(K)	RV051 2
	UDTXV(K)=MINRV(K)/TV(K)	RV052
	DO, 9972, I=1, MXMV(K)	RV052 1
	CXV(S)*DER1(QNVAR(K, I))=PCOV(K, I)*SNVAR(K, I)*UDTXV(K)	RV053
	GXV(S)*VPARI(J)*DER1(QPVAR(K, I))=PCOV(K, I)*SPVAR(K, I)*UDTXV(K)	RV054
	GXV(S)*VDISP(J)*DER1(QDVAR(K, I))=PCOV(K, I)*SDVAR(K, I)*UDTXV(K)	RV055
	CXV(S)*SNVAR(K, I)=QNVAR(K, I+1)/TVMIN(K)-VPARI(I)*VPARI(J)/TVMIN(K)	RV056
	1*GQVAR(K)+VDISP(I)*VDISP(J)/TVMIN(K)*FQVAR(K)+VDISP(I)*VPARI(J)/TVMIN(K)*FQVAR(K)-VPARI(I)*VPARI(J)/TVMIN(K)*FQVAR(K)	RV057
	2*GQVAR(K)*FQVAR(K)-VPARI(I)*VPARI(J)/TVMIN(K)*FQVAR(K)+VDISP(I)*VPARI(J)/TVMIN(K)*FQVAR(K)	RV058
	GXV(S)*VPARI(J)*SPVAR(K, I)=QPVAR(K, I+1)/TVMIN(K)-VPARI(I)/TVMIN(K)	RV059
	1*QDVAR(K)+VDISP(I)/TVMIN(K)*SDVAR(K)	RV060
9972	GXV(S)*VDISP(J)*SDVAR(K, I)=QDVAR(K, I+1)/TVMIN(K)+VDISP(I)/TVMIN(K)	RV061
	1*FQVAR(K)-VPARI(I)/TVMIN(K)*GQVAR(K)	RV062
	CXV(S)*DER1(QNVAR(K, J))=-VPARI(J)*PCOV(K, J)/TVMIN(K)*GQVAR(K) *UD	RV063
	1TXV(K)+VDISP(J)*PCOV(K, J)/TVMIN(K)*FQVAR(K)*UDTXV(K)	RV064
	GXV(S)*VPARI(J)*DER1(QPVAR(K, J))=-PCOV(K, J)/TVMIN(K)*GQVAR(K) *UD	RV065
	1TXV(K)	RV066
9973	GXV(S)*VDISP(J)*DER1(QDVAR(K, J))=+PCOV(K, J)/TVMIN(K)*FQVAR(K)*UDTXV(K)	RV067
	1V(K)	RV068
		RV069
		RV070
		RV071

Fig. 5 Listing of the APACHE subroutine for variable time delays.



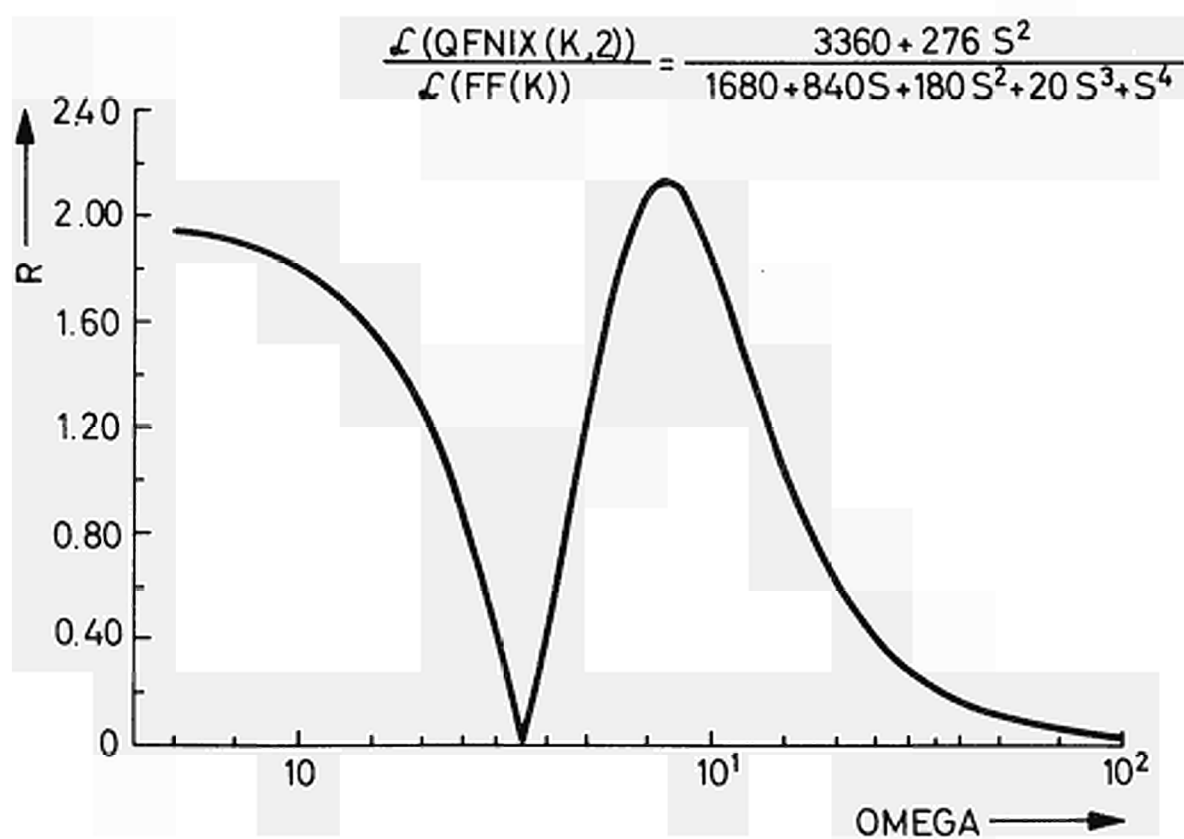
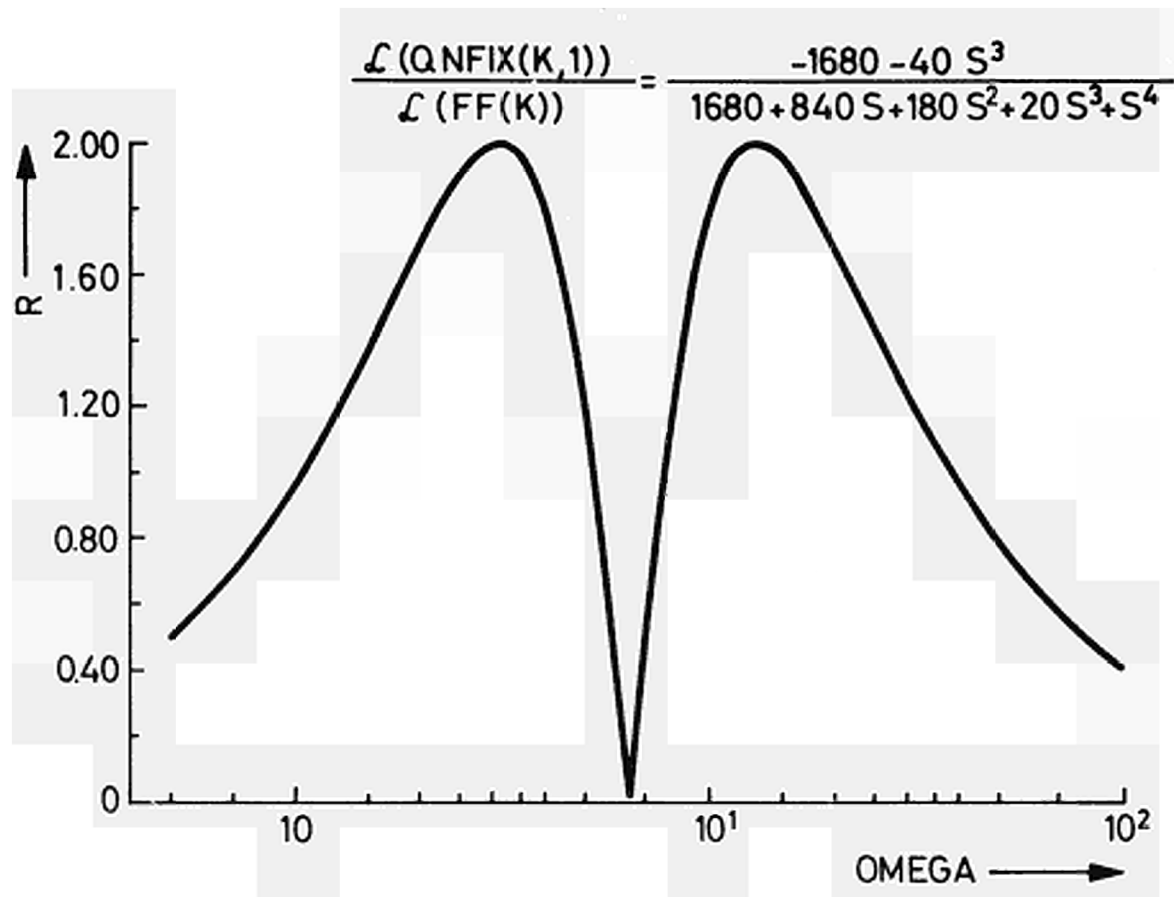


Fig. 6 Amplitude vs frequency diagram of  $\frac{\text{QNFIX}(K,1)}{\text{FF}(K)}$  and  $\frac{\text{QNFIX}(K,2)}{\text{FF}(K)}$ .

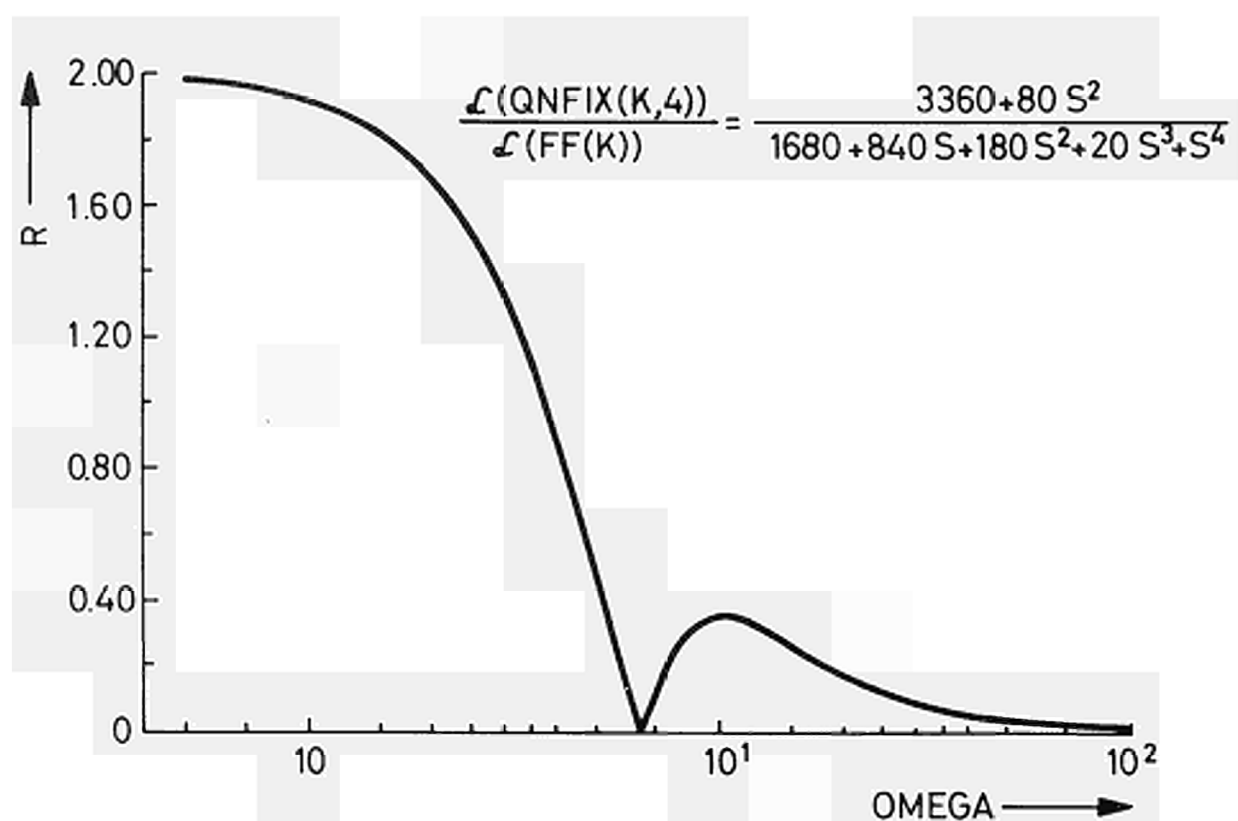
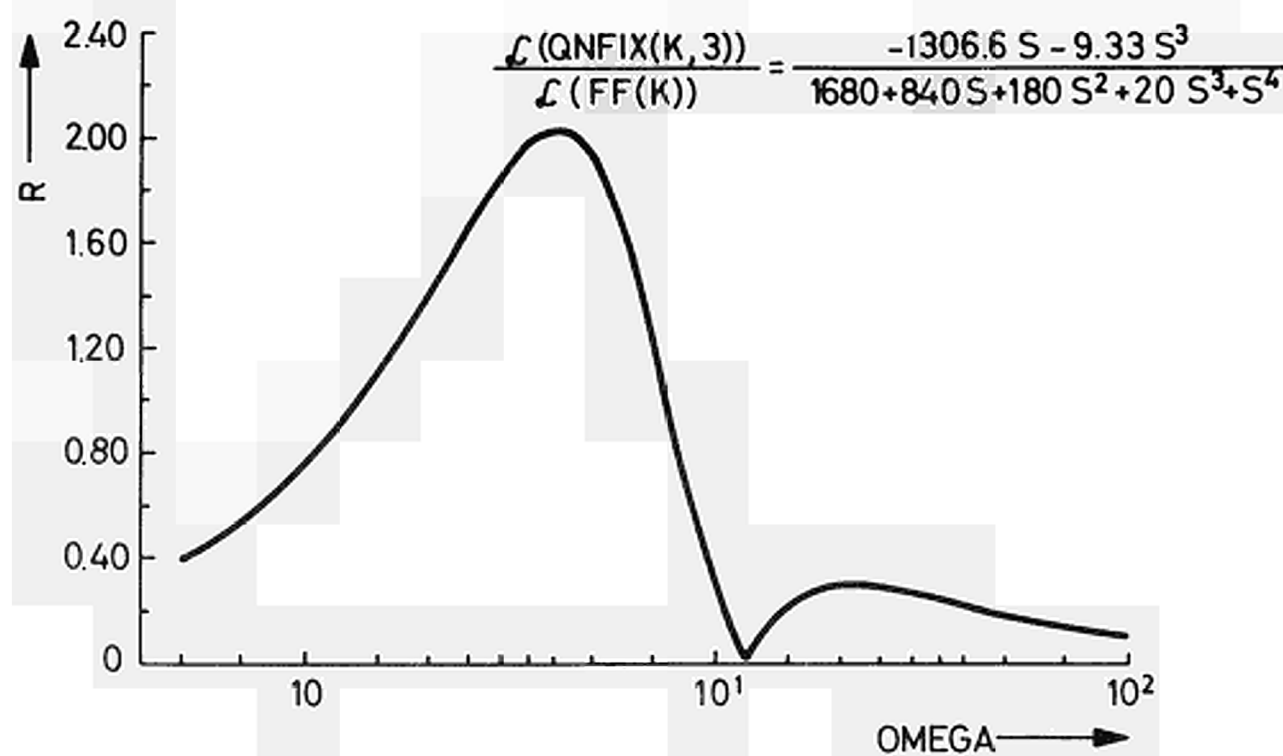


Fig. 7 Amplitude vs frequency diagram of  $\frac{\text{QNFIX}(K,3)}{\text{FF}(K)}$  and  $\frac{\text{QNFIX}(K,4)}{\text{FF}(K)}$ .

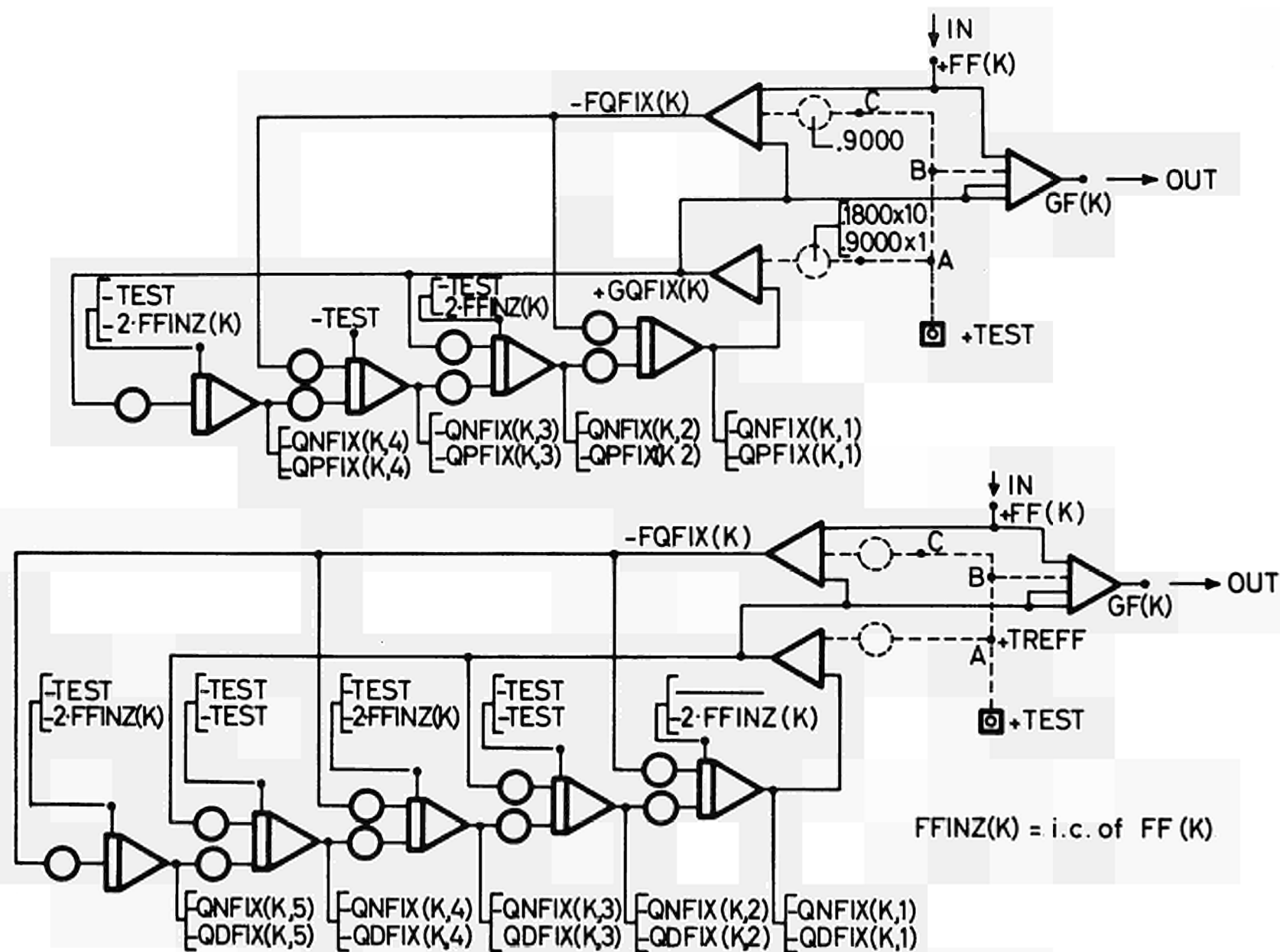


Fig. 8 Fixed time delay Padé circuits as deduced from the APACHE output listing.

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(To be published)



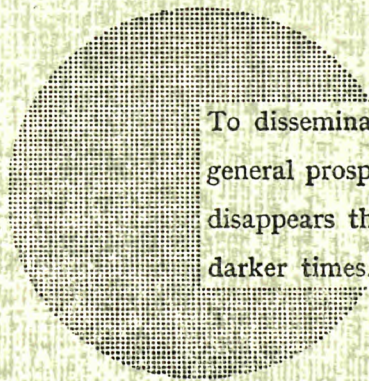
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Alfred Nobel



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